



Sensor Networks for Measurement and Data Reporting

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Motivation

ü Job request from various sources

- | Sensor generated data
- | Data intensive computing
- | Computational grids

ü Challenge

- | Minimum processing time
- | Minimum energy use

ü Solution

- | Efficient job scheduling strategies
 - r Divisible Load Theory

Topics

- q Research Purpose
- q Sensor Networks
- q Divisible Load Theory
- q Scheduling Strategies
- q Performance
- q Conclusion

Research Purpose

ü Optimal Job Scheduling

- | Measurement time
- | Computation time
- | Communication time

ü Modeling

- | Divisible load theory
- | Markov chain models

ü Application areas

- | Sensor networks
- | Grid networks

Sensor Networks

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Sensor Networks

Sense \Leftrightarrow Process \Leftrightarrow Communicate

ü Integrated collaboration

- Gathering accurate information in a distributed manner from
 - r Inaccessible geographic area
 - r Disaster area
 - r Industrial location
- Perform task
 - r Data gathering
 - r Event detection and estimation
- Relay Information
 - Network topologies
 - Access technology

Sensor Networks (cont)

ü Implications

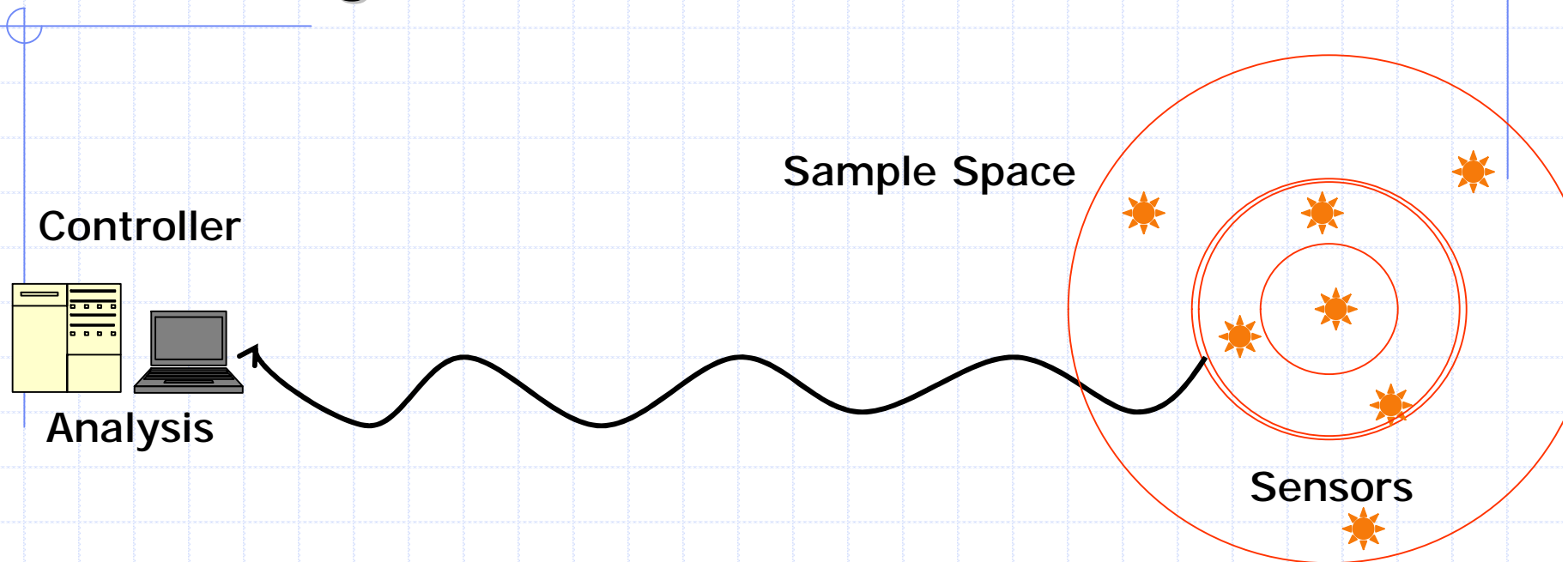
- Fixed wire-less network
- Avoid long distance communications
- No user attendance
- Deployed in large numbers
- Requires simple energy efficient routing

ü Objectives

- | Minimum processing time
 - r Fast response
- | Energy use optimization
 - r Extended life time

Our Scope

üSensing for measurement methods



- Sensors distributed on specific sample space
- Sensors measure stream of data (frequency)
- sensors report to the controller
- Sensors without computing power
- Sensors are power and energy limited

Our Scope

ü Measurement Applications

I Military Application

- r Space with large frequency range
- r Each sensor measures specific frequency and directional range

I Very Large Array of Antennas (VLA)

- r Produce image of radio sky
 - r Wide range of frequencies and directions
- r Extra-terrestrial signals

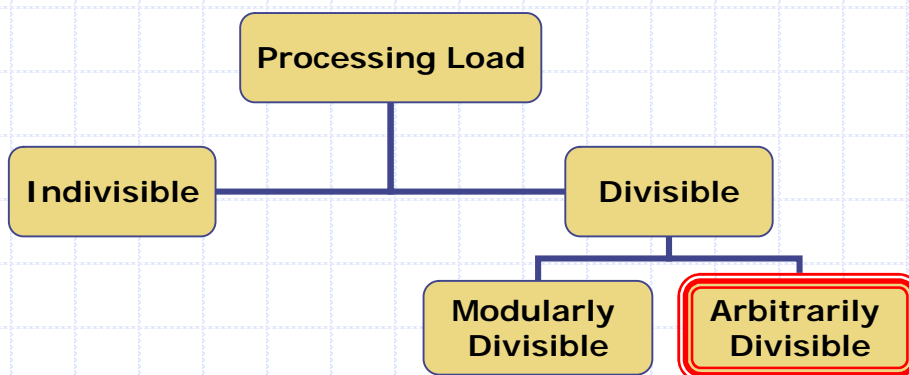
Divisible Load Theory

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Divisible Load Theory

- ü Process of arbitrarily partitioning a large load
 - | Shortest possible completion time
 - | Lowest energy use

- ü Need for efficient partitioning
 - Prevalence of multiple processor systems
 - Data intensive computing
 - Ubiquity of sensor generated data



- r Signal processing
- r Image processing
- r Sensor & Experimental data processing

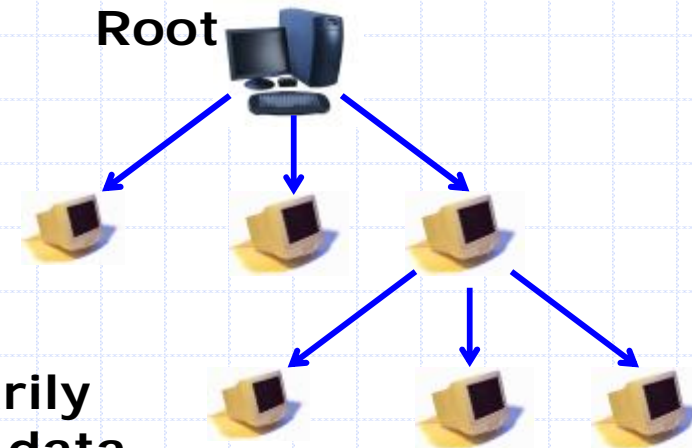
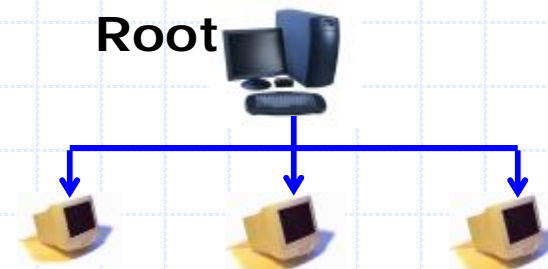
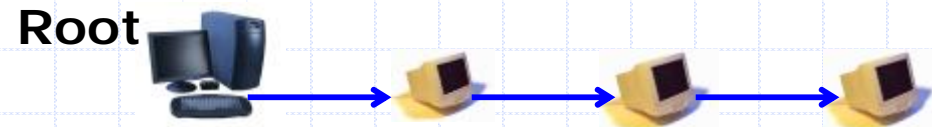
System Model

ü Topologies

- | Linear daisy chains
- | Bus networks
- | Tree Networks

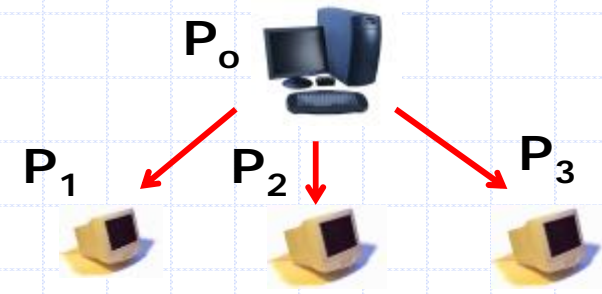
ü Assumptions

- | Load arrives at one particular processor (Root Processor)
- | Load can be partitioned arbitrarily
- | No precedence relation among data
- | Total load is normalized to be one



System Model

DLT Model

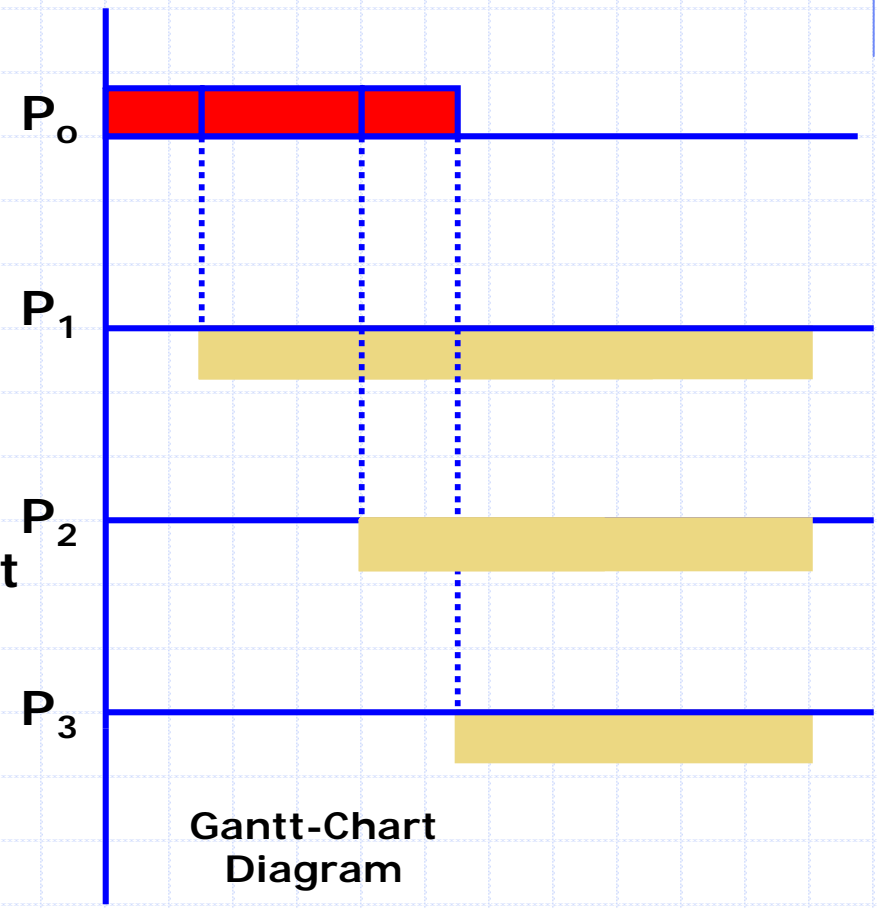


ü Types of Loads

- | Communication loads
- | Computation loads
- | Measurement loads

ü Scheduling policies

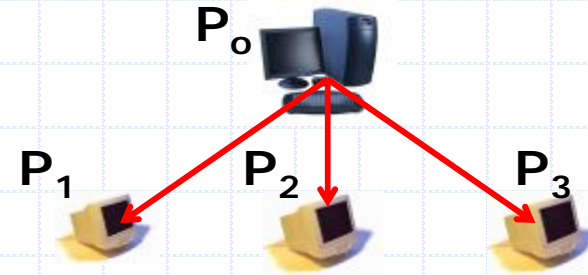
- ü Sequential
- | Concurrent/Simultaneous
- | Single or multi-installment



Gantt-Chart Diagram

System Model

DLT Model

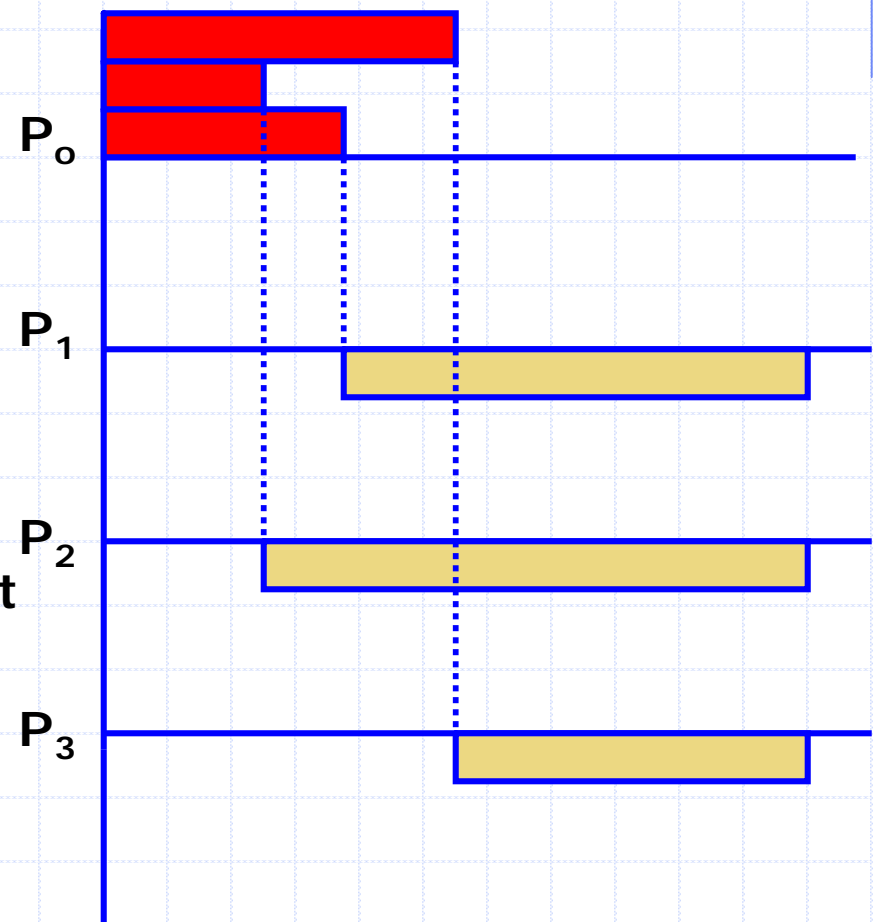


ü Types of Loads

- | Communication loads
- | Computation loads
- | Measurement loads

ü Scheduling policies

- | Sequential
- ü Concurrent/Simultaneous
- | Single or multi-installment



System Model

ü Optimality

- | Defined in the context of
 - r Specific interconnection topology
 - r Load scheduling policies

ü Solution

- | Force **ALL** processors to finish at the same time.

Notations and Definitions

- a_i – fraction of load processed by processor i
- w_i – inverse computing speed of the i^{th} processor
- y_i – inverse measurement speed of the i^{th} processor
- z_i – inverse transmission speed of the i^{th} link
- T_{cp} – computation intensity
- T_{cm} – communication intensity
- T_{ms} – measurement intensity
- $a_i w_i T_{cp}$ – time to process the i^{th} load fraction on the i^{th} processor
- $a_i z_i T_{cm}$ – time to transmit the i^{th} load fraction on the i^{th} link
- $a_i y_i T_{ms}$ – time to measure the i^{th} load fraction on the i^{th} processor
- T_i – total time elapsed by the processor i
- T_f – finish time of the process

Scheduling Strategies

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Data measurement and reporting strategies

ü Three types

- | Simultaneous Measurement Start, Sequential Reporting
r (SMS²R)
- | Simultaneous Measurement Start, Simultaneous Reporting Time
r (SMS²RT)
- | Concurrent Measurement and Reporting
r (CMR)

SMS²R

ü Single level tree network

I Three phases

r Job assignment phase

- r Each processor will receive its measurement instruction

r Measurement phase

- r Each processor starts measurement

- r Measurement time è $a_i y_i Tms$

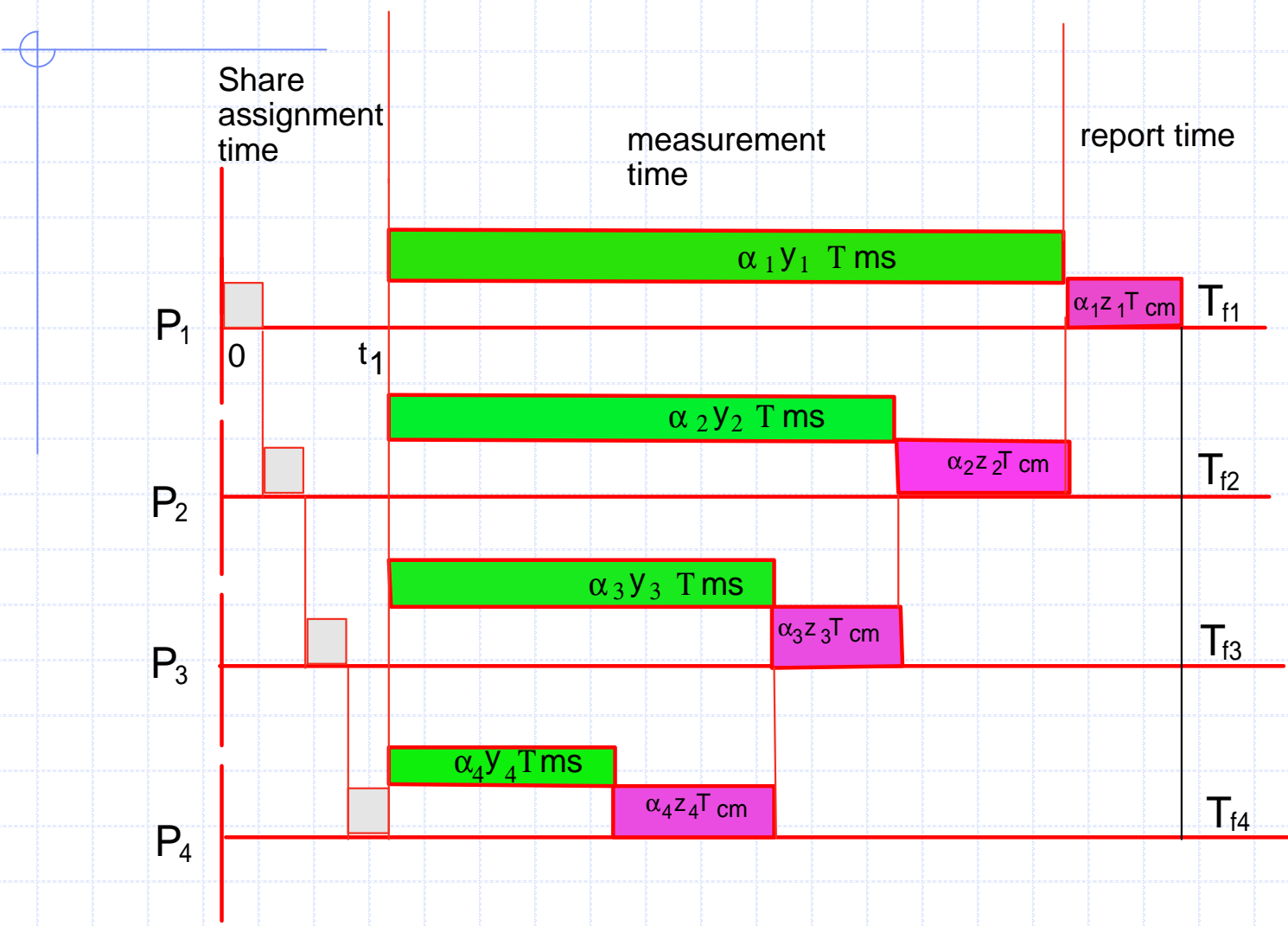
r Data Reporting phase

- r Each processor starts reporting measured data

- r Reporting is done sequentially

- r Reporting time è $a_i z_i Tcm$

SMS-SqR Timing Diagram



SMS²R

ü Set of equations for the finish time

- | $P_1 \Rightarrow T_{f1} = a_1 y_1 T_{ms} + a_1 z_1 T_{cm}$
- | $P_2 \Rightarrow T_{f2} = a_2 y_2 T_{ms} + a_2 z_2 T_{cm}$
- | $P_3 \Rightarrow T_{f3} = a_3 y_3 T_{ms} + a_3 z_3 T_{cm}$
- | $P_4 \Rightarrow T_{f4} = a_4 y_4 T_{ms} + a_4 z_4 T_{cm}$

ü Normalization equation

- | $a_1 + a_2 + a_2 + a_4 = 1$

ü Different finish times

- | $T_{f(i+1)} = a_i y_i T_{ms}$

ü Minimum finish time T_f

$$T_f = t_1 + \frac{y_1 T_{ms} + z_1 T_{cm}}{1 + \sum_{i=2}^N \prod_{j=2}^i s_j}$$

$$s_j = \frac{y_{j-1} T_{ms}}{y_j T_{ms} + z_j T_{cm}}$$

SMS²RT

ü Single level tree network

I Job assignment phase

- r Each processor will receive its measurement instruction

I Measurement phase

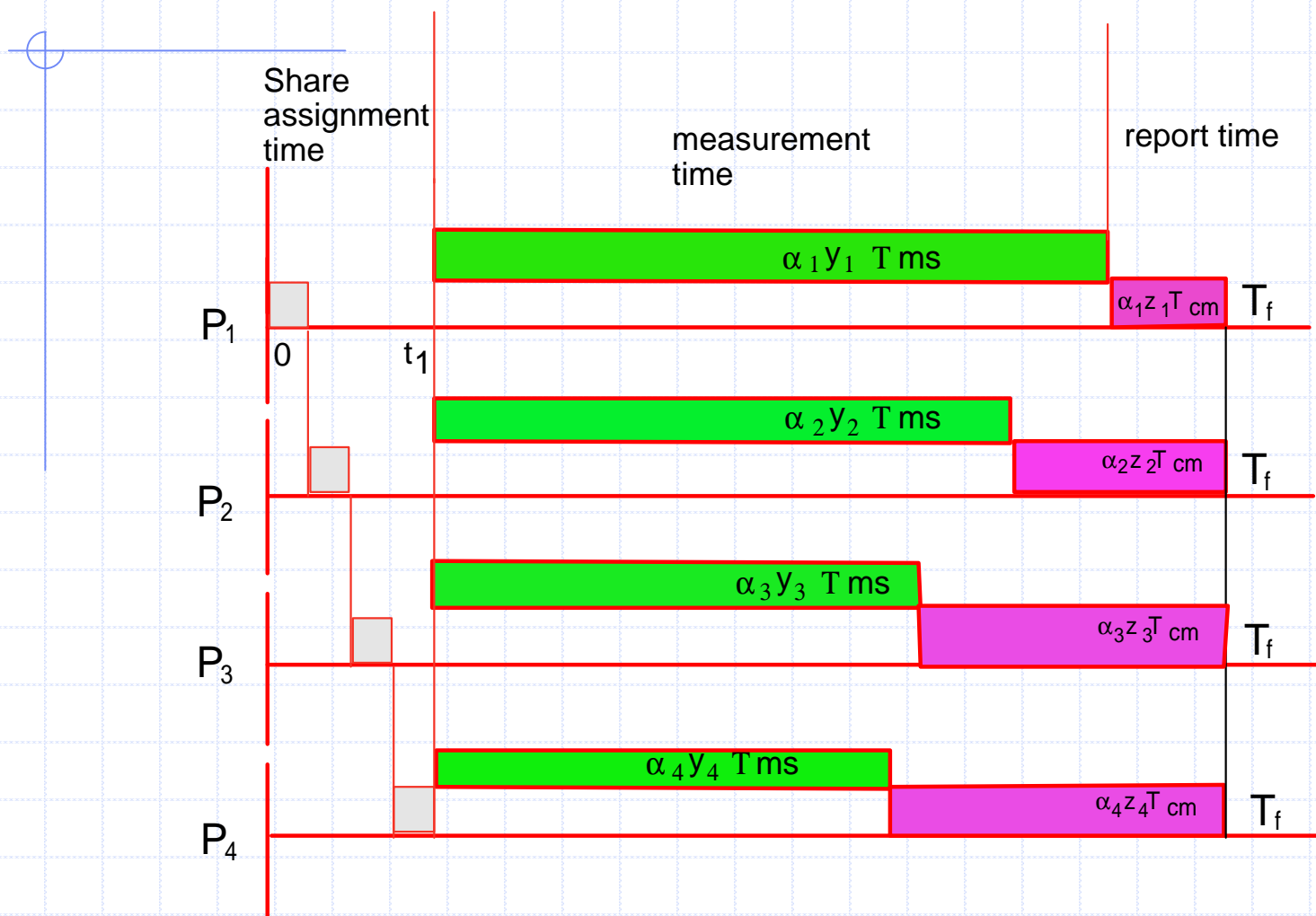
- r Each processor starts measurement

I Data Reporting phase

- r Each processor starts reporting measured data
 - r Simultaneously
 - r Same finish time

ü Each node has separate channel

SMS-SRT Timing Diagram



SMS²RT

ü Set of equations for the finish time

| $P_1 \Rightarrow T_1 = \alpha_1 y_1 T_{ms} + \alpha_1 z_1 T_{cm}$

| $P_2 \Rightarrow T_2 = \alpha_2 y_2 T_{ms} + \alpha_2 z_2 T_{cm}$

| $P_3 \Rightarrow T_3 = \alpha_3 y_3 T_{ms} + \alpha_3 z_3 T_{cm}$

| $P_4 \Rightarrow T_4 = \alpha_4 y_4 T_{ms} + \alpha_4 z_4 T_{cm}$

ü In this case, we have:

| $T_1 = T_2 = T_3 = T_4 = T_f$

ü Normalization equation

| $\alpha_1 + \alpha_2 + \alpha_2 + \alpha_4 = 1$

ü Minimum finish time

$$T_{fi} = t_1 + \frac{(y_i T_{ms} + z_i T_{cm}) \left(\frac{1}{r_i}\right)}{\sum_{i=1}^N \left(\frac{1}{r_i}\right)},$$

$$r_i = y_i T_{ms} + z_i T_{cm}, i=1,2,3,..$$

CMR

ü Single level tree network

I Job assignment phase

- r Each processor will receive its measurement instruction

I Measurement phase

- r Each processor starts measurement

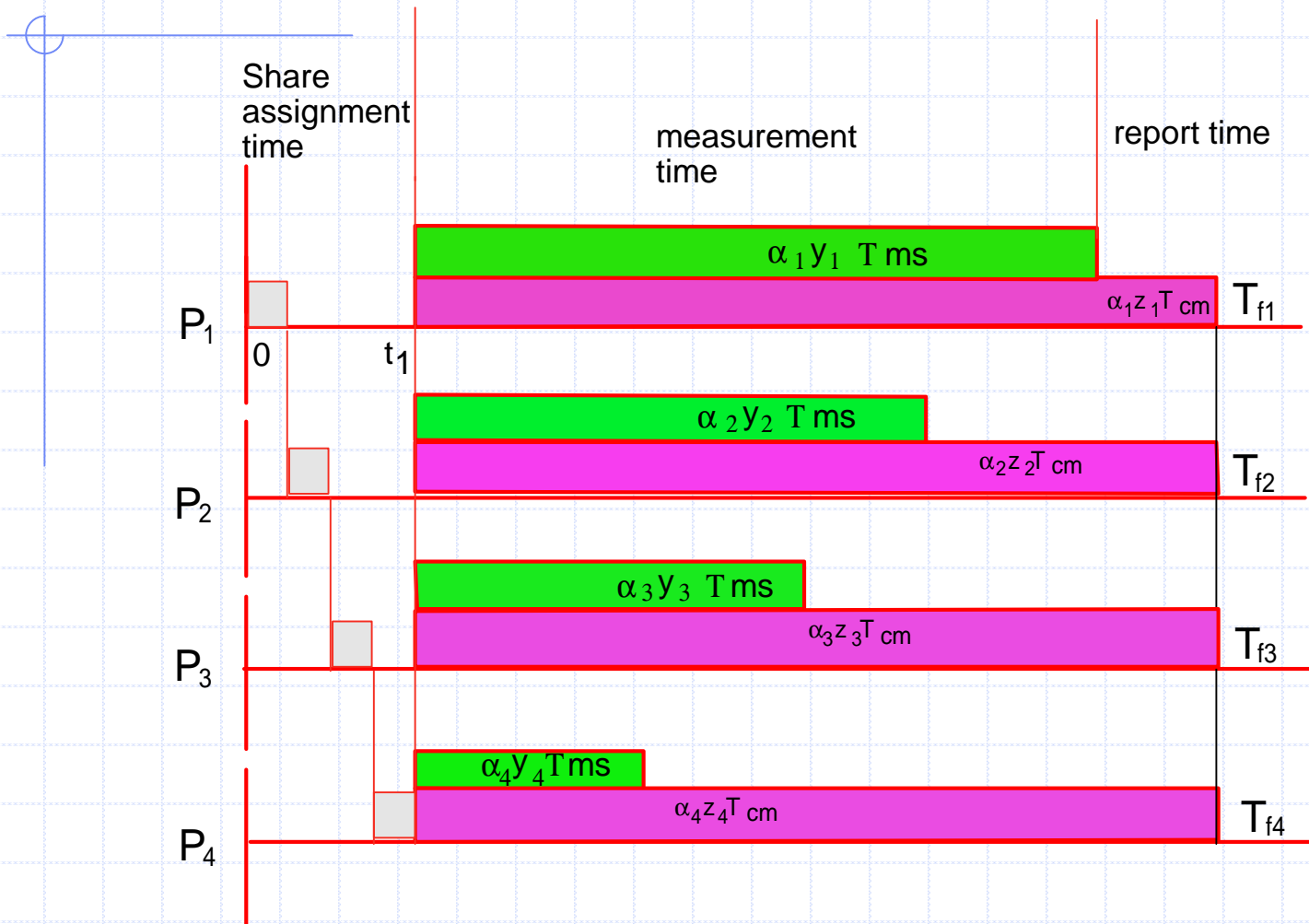
I Data Reporting phase

- r Simultaneous data reporting
- r Same finish time

ü Each node has separate channel

ü Concurrent measurement and data reporting

CMR Timing Diagram



CMR

ü Set of equations for the finish time

| $P_1 \Rightarrow T_1 = \alpha_1 z_1 T_{cm}$

| $P_2 \Rightarrow T_2 = \alpha_2 z_2 T_{cm}$

| $P_3 \Rightarrow T_3 = \alpha_3 z_3 T_{cm}$

| $P_4 \Rightarrow T_4 = \alpha_4 z_4 T_{cm}$

ü Same time finish

| $T_1 = T_2 = T_3 = T_4 = T_f$

ü Normalization equation

| $\alpha_1 + \alpha_2 + \alpha_2 + \alpha_4 = 1$

ü Minimum finish time

$$T_{fi} = t_1 + \frac{(z_i T_{cm}) \left(\frac{1}{z_i}\right)}{\sum_{i=1}^N \left(\frac{1}{z_i}\right)},$$

Performance

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Performance

ü Finish time Vs. Number of processors

ü Methodology

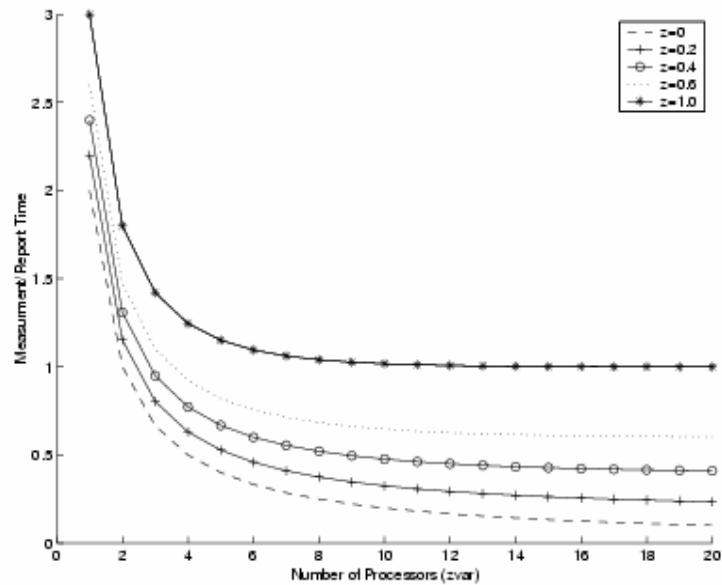
- | z_i is varied and y_i fixed
 - r Impact of link speed
- | z_i is fixed and y_i varied
 - r Impact of measurement speed

Performance

I SMS²R

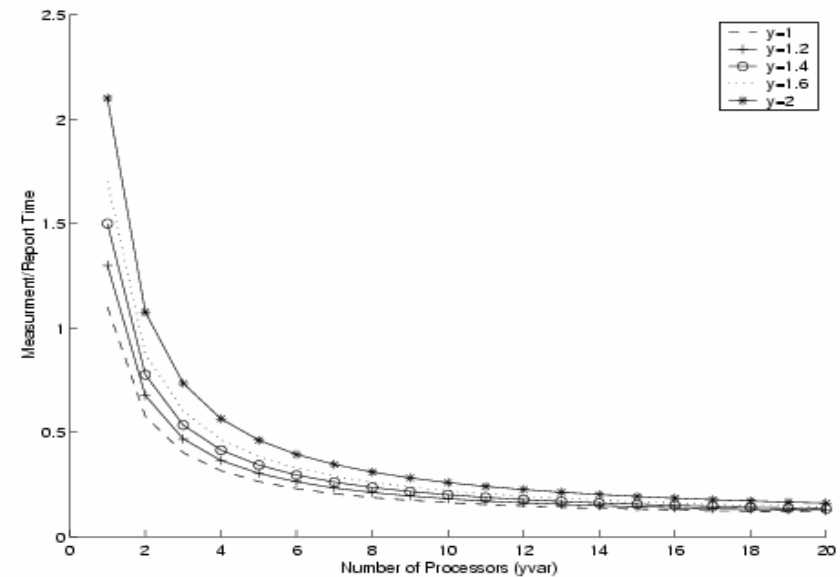
r z_i varied from 0 to 1

r y_i fixed to be 2



z_i fixed to be 0.1

y_i varied from 1 -2



Performance

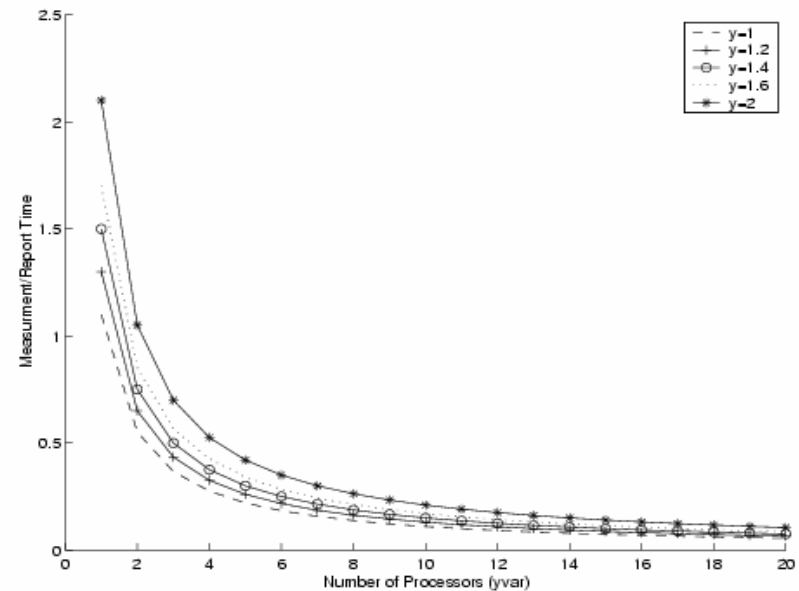
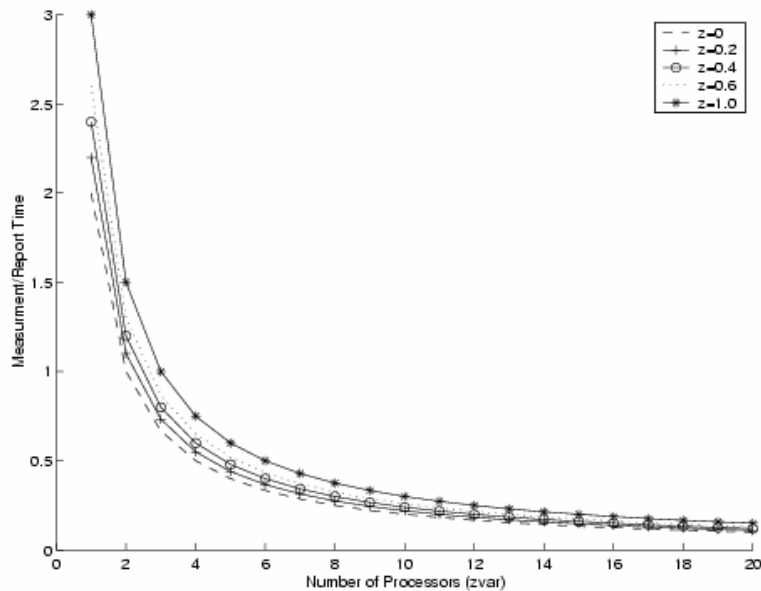
I SMS²RT

r z_i varied from 0 to 1

r y_i fixed to be 2

z_i fixed to be 0.1

y_i varied from 1 -2

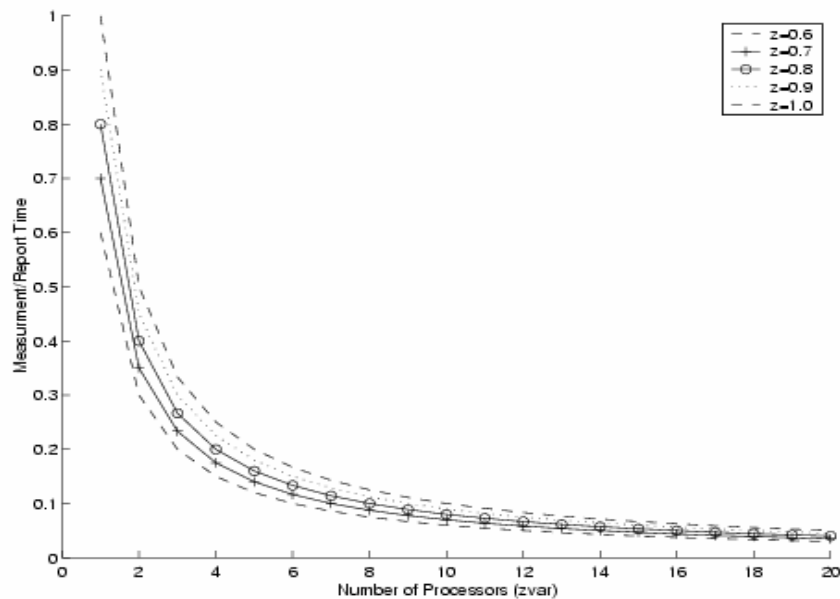


Performance

I CMR

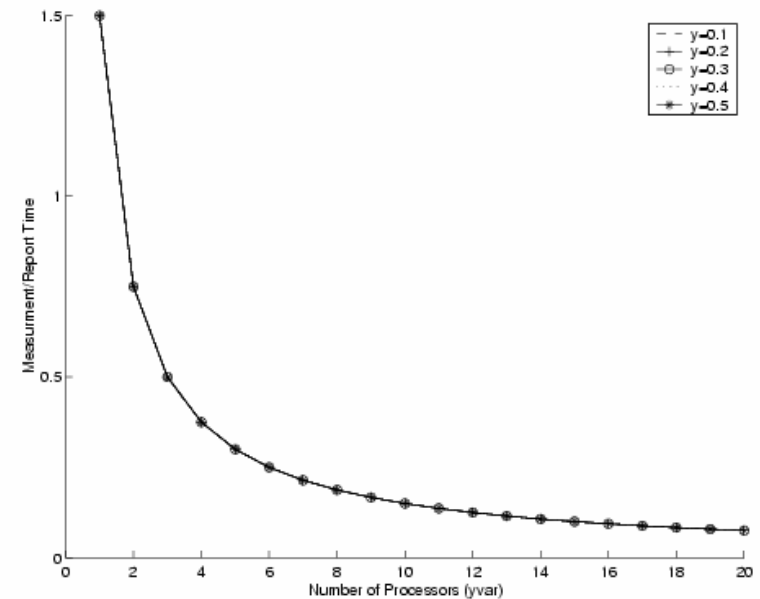
r z_i varied from 0 to 1

r y_i fixed to be 2



z_i fixed to be 0.5

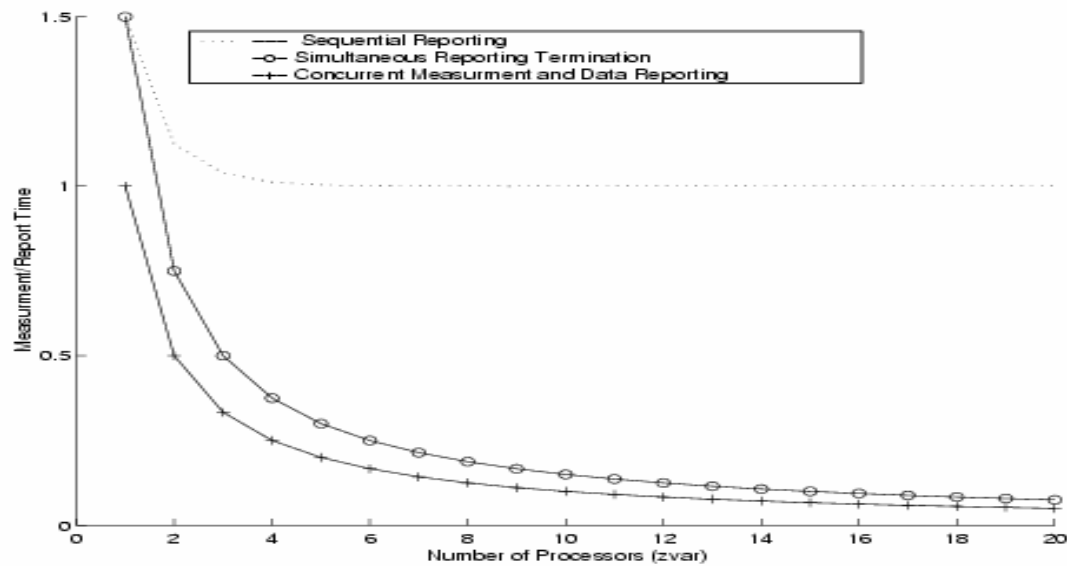
y_i varied from 0.1–0.5



Performance

comparison

CMR performs WELL!



Conclusion

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Conclusion

ü Ongoing research

- | Monetary cost based energy use
- | Fault and rescheduling
- | STAR experiment in the RHIC project
 - r Collaborative work with BNL

ü Future research

- | Experimental level research
- | BlueGene/L , IBM
 - r Massively-parallel supercomputing
- Development of Integrated Scheduling Algorithm for Distributed Mobile Sensor Network
 - r <http://www.tech.uh.edu/isgrin/index.htm>